#### MICROPOLIS USERS GROUP

#### MUG Newsletter # 15 - October 1981

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# WRITING GOOD PROGRAMS (PART 3)

#### by Burks A. Smith of DATASMITH Box 8036, Shawnee Mission, KS 66208

Last month we discussed techniques for writing code that produces self-documenting, modular programs that are easy to maintain and less likely to have bugs. I have concentrated on planning techniques and general styles of writing programs rather than language details because good style encourages good content. In order to become fluent in any language (French, Basic, or whatever) you must use it often. This is really the only way you can learn programming, but starting out with proper planning and style will make it easier.

#### SERVICING OPERATOR INPUTS

Most programs for business or scientific purposes have some part of the input data entered directly from the keyboard by a human. This input data may vary from a simple "yes" or "no" answer to large volumes of figures for actual processing. The program must be able to get along with its human users, and a very important aspect of this is the ability to tolerate human error with out "crashing" or making a mistake itself. Everything coming from the keyboard should be checked for errors, and if an error is detected the operator should be politely informed and given another chance.

The simplest case is a question which requires a "yes" or "no" answer. First of all, the question put to the operator should be easy to understand and indicate what kind of a response is expected. "DO YOU WANT TO USE THE PRINTER (Y OR N)?" is a typical example of this. In response to the above question, a string of characters is INPUT from the keyboard. It is expected that the operator entered a "Y" or an "N", since this is what was instructed, but as the programmer, you must be prepared for anything and allow a little leeway too. For example, are you going to reject anything except the single characters "Y" or "N"? That may be too rigid. What if the operator enters an "N" followed by a blank space? N-space does not equal N. Will you take "YES" or "NO" too? What if the operator just hits the RETURN key?

All the possibilities must be considered when coding even a simple question. The program's treatment of the answer depends, to a great extent, on the application. However, I prefer to use an algorithm that only considers the first character entered. Use a LEFT\$ function to isolate the first character in the string. If it is a "Y", assume "yes"; if it is an "N", assume "no". Otherwise, assume an invalid response, report to the operator that you can't understand, and ask again. Repeat the loop until you get a valid answer. In MICROPOLIS BASIC, just hitting the return key in response to an INPUT statement leaves the value of the input variable unchanged. Therefore, it is wise to initialize the input variable to "", or a default answer, before the INPUT statement or the results may not be predictable.

Operator response to a menu of numbered choices is a little more easy to filter out. The value input at the keyboard will be a number, so a numeric variable is used. BASIC won't allow anything but a number in this case, so you know at least a number will be returned by the INPUT statement. All you have to do is make sure the number entered is within the the limits of the menu choices. For really foolproof menus, you should avoid using an integer variable as the variable input. I seldom heed my own advice in this respect, but keep in mind that an integer variable is stored as a fixed number of digits, while a real (floating point) variable is stored in scientific notation. With the default SIZES, your program will crash with a CONVERSION ERROR if the operator enters 500000 or more through an INPUT statement with an integer variable. Try it.

When the operator needs to enter large amounts of data, it is absolutely necessary that the operator be given an opportunity to personally verify the data before it is stored away or used in a calculation. The easiest way to do this is to re-display it after it is input and ask the operator if it is OK. Facilities should be provided to change a single item, without having to re-enter all the data, if an error is detected. Programs that use a database on diskette should have facilities to manually update ANYTHING on the diskette, and to manually recover from errors made when erroneous data is used in calculations. The excuse, "It's in the computer and we can't change it," is a sure sign of a bad program. In fact, it should be easier to change data in a computer than data on paper. Computers are supposed to make things easier, not more difficult!

#### SERVICING DISK ACCESSES

Besides errors from data entered in INPUT statements, the most common cause of programs crashing can be attributed to disk errors. Not real I/O errors due to hardware, but things people do or forget to do. The most common ones are DRIVE NOT UP and FILE NOT FOUND, caused by the program expecting to find something that has not been provided by the operator, but there are others. See the ERR function in the MICROPOLIS BASIC manual (table 5.5). All of these disk errors cause the program to be stopped in its tracks, and can be catastrophic to a newly updated file that is open at worst or annoying at best.

BASIC comes to the rescue with the ERROR clause in the OPEN statement for a file. <ERROR linenumber> refers to a line number to GOTO if an error is encountered in the program. The line number should contain code that determines which error has occured (the ERR function) and appropriate error handling routines. In the case of many errors, all that is needed is to inform the operator of the error and wait until the operator says its OK to retry the operation. Other errors which can't be corrected by the operator need to terminate the program in an orderly way by closing files and exiting to a menu, etc. For example, a DISK FULL error handler definitely needs to close the file it was updating so no data will be lost. Beyond that, it depends on the program and the application. Some programs might instruct the operator to mount a fresh diskette and pick-up where they left off, while others would have to be terminated.

Trying to imagine all the errors that could possibly occur and devising software recourse is not easy. Murphy's Law states: "If something can go wrong, it will go wrong," and I read somewhere that: "It is impossible to make anything foolproof because fools are so ingenious." Good error handling routines take programming time that could be used for more exicting tasks, but chances are you will eventually need it. Also, no matter how hard you try, you will probably overlook something. Even so, don't be discouraged; any error handling is better than none and your programs will be more reliable as a result.

NEXT MONTH we'll cover some nitty-gritty details of how to design loop structures.

## TABLE HANDLING TECHNIQUES - Part 2

#### by N. P. Dembinski 8618 Essex Ave., Chicago IL 60617

Last month we summarized some of the programming techniques and guidelines useful for efficient table handling. This month, a description of each technique is presented. Also, the flowchart on the last page is intended to aid programmers in the selection of the most suitable table handling techniques for the specific problems and applications.

#### A. DIRECT TECHNIQUE

The table contains as many entries as the value of the highest numeric table code being processed. The code being searched for literally becomes the index for accessing the corresponding table entry.

### LIMITATIONS

The codes being processed must contain numeric digits. The codes usually are not in straight numeric sequence, (e.g., 001, 002, 003, etc.). Gaps occur between codes and the table size becomes prohibitive. Input data must be read at the beginning of each step to fill the table.

#### B. SERIAL TECHNIQUE

Search the table beginning with the first position, for each transaction and incrementing the index by one, or coding repetitive compares. Entries do not have to be in sequence but may be ordered by testing for the most frequently accessed items first.

#### LIMITATION

When the table size exceeds 30 entries, the execution time to test all the entries becomes prohibitive, and the input data must be read at the beginning of each step to fill the table.

## C. BINARY TECHNIQUE

The binary search technique is a good choice for over 60 entries if memory is critical. The argument to be found is first compared against the middle of the table. On the basis of this comparison, the upper or lower half of the table will be compared against. The table will be halved until an equal condition occurs or the size of the table is reduced to zero. The table arguments must be in sequence. A count of the number of entries (easily generated when the table is filled) is also required. The table size should be one less than the power of two.

#### D. INCREMENT (PARTIONED) SCAN TECHNIQUE

This is most efficient method, time-wise, for over 60 entries. Search increment entries must be in sequence, but entries within a group do not have to be in order. When searching a table with this organization, locate the group of entries wanted and move that group of entries to a work area for a detailed serial search.

#### E. SUBSCRIPTED ITERATIVE TECHNIQUE

This is the most common method of searching a table and is the most inefficient.

## F. LITERAL TECHNIQUE

A table does not appear in the data area, but the codes being searched for appear as literals in the procedure division of the program. A conditional statement must be written for each code to be processed. When the literal being tested equals the transaction code being processed, the corresponding data appears as a literal and is moved to a common storage area.

#### LIMITATIONS

Limited to fixed codes not subject to change since literal change requires source language changes, recompilation, and far more coding is required than in the other techniques.

With the availability of high level languages like Basic, Fortran IV, Cobol, Pascal, and a few others that may come along in the future, or some macro oriented assembler language, the literal technique becomes easy to use. The examples shown in the following text will demonstrate the use of the

Basic language. Basic is the most universal language for micro-computers that is available. These techniques have already been used or can be translated to some other high level language easily.

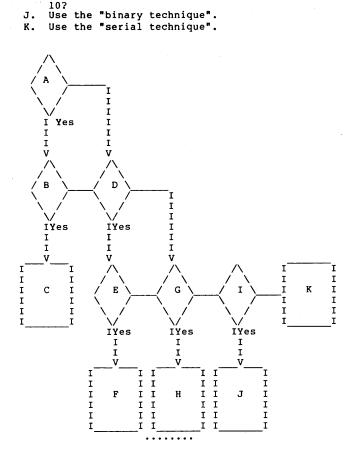
Presenting a strong case for using the literal technique, assume that a company has offices or plants in 16 cities. To conserve space on the master record, the city and state are not stored but identified by the five digit zip code number The direct approach is prohibitive since the table would require 999,999 entries allocated for 16 entries only. The amount of coding required for the literal approach is not prohibitive, there is little likelihood that the relationship of the zip code to city will ever change and the literal processing is four to five times as fast than either the serial or binary techniques. The Basic coding which would implement the literal technique follows below.

BASIC EXECUTIVE

(ETC.) GOSUB 6000 (ETC.) 6000 IGET-ADDRESS 6010 IF Z\$>"27601"THEN GOTO 6070 6020 IF Z\$ > "14205" THEN GOTO 6120 6030 IF Z\$ > "02109" THEN GOTO 6160 6040 IF Z\$ = "00902" THEN C\$ = "SAN JUAN PR" : RETURN 6050 IF Z\$ = "02109" THEN C\$ = "BOSTON, MA" :RETURN 6060 GOTO 6300 6070 IF 2\$ > "55401" THEN GOTO 6280 6080 IF 2\$ > "43216" THEN GOTO 6220 6090 IF 2\$ = "33101" THEN C\$ = "MIAMI FL" :RETURN 6100 IF Z\$ = "43216" THEN C\$ = "COLUMBUS OH" :RETURN 6110 GOTO 6300 6120 IF Z\$ > "19104" THEN GOTO 6250 6130 IF Z\$ = "15219" THEN C\$ = "PITTS. PA" :RETURN 6140 IF Z\$ = "19004" THEN C\$ = "PHIL. PA" :RETURN 6150 GOTO 6300 6160 IF Z\$ = "07102" THEN C\$ = "NEWARK NJ" :RETURN 6170 IF Z\$ = "14205" THEN C\$ = "BUFFALO NY" :RETURN 6180 GOTO 6300 6190 IF Z\$ > "85026" THEN GOTO 6280 6200 IF Z\$ = "77002" THEN C\$ = "HOUSTON TX" :RETURN 6210 GOTO 6300 6220 IF Z\$ = "57302" THEN C\$ = "MUNCIE IN" :RETURN 6230 IF Z\$ = "55401" THEN C\$ = "BLOOM IN" :RETURN 6240 GOTO 6300 6250 IF Z\$ = "21233" THEN C\$ = "BALT. MD" :RETURN 6260 IF Z\$ = "27601" THEN C\$ = "RALEIGH NC" :RETURN 6270 GOTO 6300 6280 IF Z\$ = "94086" THEN C\$ = "SUNNYVALE CA" :RETURN 6290 IF Z\$ = "96813" THEN C\$ = "HONOLULU HI" :RETURN 6300 C\$ = "" : RETURN ! BAD ZIP CODE

GUIDE FOR SELECTING TABLE LOOK-UP TECHNIQUES

- Are the keys (item no., etc.) all numeric Α. (0 thru 9)?
- Are most of the integers within the range of в. the control key required? Use the "direct technique".
- Is programming development time available to D. code the literals?
- Ε. Will the same key always relate to the same entry?
- Use the "Literal technique". F.
- G. Is it impossible to sequence the table-fill data records?
- Use the "serial technique". н.
- Does the total number of table entries exceed Ι.



#### PROGRAM DOCUMENTATION

#### by Joel Shapiro of BONJOEL ENTERPRISES Box 2180, Des Plaines, IL 60018

Most of us will spend many hours writing programs for others and our own use. However, these programs may be used so infrequently that we forget how to work with them. We'll use a lot of prompting within the program, but the overall purpose, or operation of the program may be lost because enough detail isn't given at the time they're in use.

Including a set of instructions within the program is an excellent idea and, if instructions are required by the operator, a menu selection can provide them. This takes up valuable memory and, if the program is large, you may not have enough for the program itself. Text does take up a lot of space and becomes troublesome to enter in DATA statements when you can't see the true length of each line on the screen.

Alternative to this is the use of a data file and a program or routine that will display the instructions retained in the data file. That is the approach offered here. Two programs; INSTR.GEN and INSTR.READ respectively generate and read the instruction data file.

An important feature is that the system will allow you to select the starting record of the instruction file for display and this can be keyed to the calling program. It will also allow you to return to the correct program after the instructions have been displayed. The technique for doing this is detailed within the INSTR.READ program listing. This is a good feature of the system in that it allows you to reuse the general purpose instructions you may have in your file for many different programs and routines. Also, since you may select the specific routine you wish to return to after use of the INSTR.READ program, redundant instructions are possible without duplication in the file. INSTR.GEN is charged with writing the file with the instructions as entered. It stores three (3) lines within one (1) physical file record without packing. No editing utility is provided so each line should be checked before RETURN is pressed. After three (3) lines of instructions are entered the system will pause while the record is written to the file. This process continues until you're done and a backslash is entered at the beginning of a new line to close the file. If a file already exists, additional instructions can be added to the end of the file.

Note line 1000 on the listing and you'll see how I got a modified version of Buzz's INKEY program into RAM. Of particular importance is the input subroutine in lines 20 thru 60. BASIC'S INPUT statement prints a question mark and a space at the beginning of the input line, therefore, the resultant line on the screen is offset by two spaces and won't give you an accurate indication of how the printed line will appear. This routine will but must be located in the lowest line numbers of the program to be effective. Located at the back of the program the routine can't keep up with a fast typist and characters will be missed. This is a very useful routine as placing the escape sequence for direct cursor positioning in line 20 will allow you to input the line anywhere on the screen. This allows you to draw a picture of a form and fill in the blanks directly. Like I said, it's a very useful routine.

INSTR.READ can be used as a program or changed and incorporated into a program as a routine. The operation is very simple and the technique for returning to the correct program is explained in the listing.

Note the string delimiter is changed to Ascii 255 in both programs to allow you the use of commas in the text. Line 280 in INSTR.READ provides for the return of the delimiter to the default value (a comma).

Be sure to revise the variables for SCREEN WIDTH and SCREEN HEIGHT in both programs as these values are used for proper formatting.

I hope this may encourage some of you to provide instructions as part of your program development; it certainly helps to prevent mistakes when you've forgotten how the program is supposed to work. And, as editing and a few other features were left out of these, I hope you'll keep the MUG advised if you develop any enhancements.

#### Title: INSTR.GEN

10		GOTO 410
20	>*	PRINT:B\$=""
30		C\$=FAA(1):IF ASC(C\$)=3 STOP
40		IF ASC( $C$ \$)=13 RETURN
50		IF ASC( $C$ \$)=127 OR ASC( $C$ \$)=8 THEN B\$=LEFT\$
		(B\$, LEN(B\$) - 1):GOTO'30
60		B\$=B\$+C\$:GOTO 30
70		INPUT ROUTINE IS ABOVE - MUST BE AT B
		EGINNING OF PROGRAM
80		CLEAR SCREEN
	><	PRINT CHAR\$ (26): PRINT REPEAT\$ (CHAR\$ (13)+C
		HAR\$ $(10)$ , INT $(H^*.5) - 2$ ) : RETURN
100		! REVERSE VIDEO
	><	PRINT CHAR\$ (18); :RETURN
120		! VIDEO RESET
130	><	PRINT CHAR\$(17);:RETURN
140	><	INPUT "Press RETURN When Ready ";A\$:RETUR
		N
150		PROGRAM START
160		GOSUB 90:PRINT TAB(T); "CREATE NEW INSTRU
		CTION FILE":PRINT
170		CREATES NEW FILE
180		PRINT TAB(T)"; Enter Name for New File";:
		INPUT N\$:PRINT:PRINT TAB(T); "Enter Drive
		Number for File";:INPUT A:IF A<0 OR A>3 G
		OTO 180
190		OPEN FILE
200		N\$=FMT(A, "N:9:")+N\$:OPEN 1 N\$ ERROR 210:G
200		THE CASE AND A CONTRACT OF EN I NO ERROR 210:G

00 N\$=FMT(A,"N:9:")+N\$:OPEN 1 N\$ ERROR 210:G OTO 220

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210	> PRINT:PRINT TAB(T);:GOSUB 110:PRINT "DISK	1000 >
	ERROR - ";ERR\$:GOSUB 130:PRINT TAB(T); "C	
220	<pre>orrect and ";:GOSUB 140:GOTO 180 &gt; S=RECPUT(1):GOSUB 90:PRINT TAB(T); "Enter</pre>	
220	text, line for line, exactly as":PRINT TA	1005 <
	B(T); "You wish it to appear on the scree	
230	n.":PRINT PRINT TAB(T); "Enter <cr> at the end of e</cr>	1010 1020
200	ach line.":PRINT:PRINT TAB(T); "Enter \ w	
	hen done with instructions.":PRINT:PRINT TAB(T);:GOSUB 140:GOSUB 90	Title:
240	! GET INPUTS - GETS 3 LINES AND WRITES	
050	TO FILE	10 20
250	! LINE IS COMPOSED IN B\$ - TRANSFERS TO G\$(X) ARRAY	30
260	> FOR I=0 TO 2:GOSUB 20:IF I=0 AND B\$=CHAR\$	40
270	(92) GOTO 330 IF B\$=CHAR\$(92) GOTO 310	50
280	IF LEN(B\$)>W% THEN I=I-1:B\$="":PRINT TAB(	60
	T);:GOSUB 110:PRINT "LINE TOO LONG - REEN TER":GOSUB 130:NEXT I	70 80
290	G (I)=B\$:NEXT I	90
300 310	! WRITE TO FILE > A\$=G\$(0)+Y\$+G\$(1)+Y\$+G\$(2)+Y\$:PUT 1 A\$:IF	100 110
510	$B_{3}=G_{3}(0)+13+G_{3}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)+13+G_{3}(2)+13+FO_{1}(1)$	120
	":G\$(2)="":GOTO 260	130 140
320 330	<pre>! ENDING FOR THIS ROUTINE &gt; F=RECPUT(1)-1:CLOSE 1:GOSUB 90:PRINT TAB(</pre>	140
	T); "Starting Record =";S:PRINT:PRINT TAB	150
	<pre>(T); "Ending Record =";F:PRINT:GOSUB 14 0:GOTO 530</pre>	160 170
340	! OPENS EXISTING FILE AND TRANSFERS TO	180
250	INPUT ROUTINE ABOVE	190
350	> GOSUB 90:PRINT TAB(T); "ADD TO EXISTING I NSTRUCTION FILE":PRINT	200
360	> PRINT TAB(T); "Enter Name for Existing In	210
	struction File";:INPUT N\$:PRINT:PRINT TAB (T); "Enter Drive Number for File";:INPUT	210
	A:IF A<0 OR A>3 GOTO 180	220
370	N\$=FMT(A,"9:")+N\$:OPEN 1 N\$ ERROR 380:GOT 0 220	230
380	> PRINT: PRINT TAB(T);:GOSUB 110: PRINT "DISK	240
	ERROR - ";ERR\$:GOSUB 130:PRINT TAB(T); "C orrect and ";:GOSUB 140:GOTO 360	250
390	> GOSUB 90:A\$="PROGRAM TERMINATED":PRINT TA	260
	B(FNA(A\$));:GOSUB 110:PRINT A\$:GOSUB 130: PRINT:PRINT:PRINT:END	270
400	PROGRAM COMES HERE FOR INITIALIZATION	280
410	<pre>&gt; DIM C\$(1),A\$(250),G\$(2,81),B\$(80):GOSUB 1 000</pre>	290 300
420	SCREEN WIDTH	310
430	W=80	320 330
440 450	₩%=₩-1 ! SCREEN HEIGHT	340
460		1000 1010
470	DEF FAA=16R4E:DEF FNA(X\$)=INT((W-LEN(X\$)) /2):Y\$=CHAR\$(255):STRING Y\$:T=INT(W*.25)	1020
480	A\$="INSTRUCTION FILE GENERATOR":GOSUB 90:	1030 1040
490	PRINT TAB(FNA(A\$));A\$:PRINT PRINT "This program generates a file cont	1040
450	aining the instructions for":PRINT "opera	1060 1065
	ting a program. The file is read by calli ng INSTR.READ":PRINT "from your program."	1005
	:PRINT	1080 1090
500	PRINT "INSTR.READ and the INSTRUCTION FIL E must reside on drive 0 for":PRINT "prop	1095
	er operation. INSTR.READ is called with a	1100 1110
	PLOADG ":PRINT "'INSTR.READ' within the c alling program. If you wish to return"	1120
510	PRINT "to the proper program after the in	1130
	structions are displayed,":PRINT "follow the remarks contained within the listing	1140 1150
	for INSTR.READ.":PRINT:PRINT TAB(T);:GOSU	1155
520	B 140 ! Main Menu	1160 1170
520	> GOSUB 90:PRINT TAB(T); "Functions Availab	1180
540	le":PRINT PRINT TAB(T); " 1) Create New Instruction	1190 1200
540	File"	1210
550	PRINT TAB(T); " 2) Add To Existing Instru	1220 1230
560	ction File" PRINT TAB(T); " 3) Terminate Program"	1240
570	PRINT:PRINT TAB(T); "Select Function You Desire";:A\$=FAA(1):IF A\$=CHAR\$(3) STOP	1250 1255
580	IF A\$<"1" OR A\$>"3" GOTO 530	1260
590	A=VAL(A\$):ON A GOTO 160,350,390	1270 1280
990	! LOADS MODIFIED INKEYS ROUTINE STARTIN G AT LOCATION 16R4E	1285
		1290 1300
		1310

00 05 10 20		1 1 1	778 RI B <b>%,</b>	32 =0 2) 1: \$(	2A3 ) T ): FO (B\$	01 0 8 8 8 98	CD 12 =B I= %, 1	7F :P 8+ 13 2) BO	07 OK 2: T ):	00 E( NE 0 B% 0E	00 A% XT 26 =B L	00 +I I :P %+	C9 )= OK 2:	": VA E( NE	A % L ( A % XT	=1 "1 +I I	6R4 6R" )=V :RE	E: +M	B9 110	CD =1: )\$ () )\$ ()	:FO A\$,	
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230 240		>		Т	1	G\$	1A (0) T 9	),(	;\$	(1)	),(	G\$	(2)	):E	PUI			\$ (	0)	:PU	т 9	)
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100			!		PO	KE	(16	5R4	D)	= 5	5 : F		DAE	G"	IN	IST	R.I	REJ	AD'			
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#### INTERRUPT HARDWARE FOR REAL TIME CLOCK

by Howard Rowland WBlAJX 79 Ivan Street Apt 65 North Providence RI 02904

## REAL TIME CLOCK INTERRUPT

This article describes the implementation of a simple real time clock on a Micropolis based system.

A brief description of how the CPU handles interrupts, then a description of hardware that can be used will be presented.

#### 8080 INTERRUPTS

When the 8080 receives an interrupt request while interrupts are enabled, it finishes the current instruction and starts an interrupt acknowledge cycle. This cycle is identified by the output of the INTA status signal on S-100 pin 96. During this cycle, the interrupt hardware in the system will place an instruction on the data in bus. Typically, this is a restart instruction, which is a one byte call to a set address in low memory. After receiving this instruction, the processor will push the contents of the program counter onto the stack. A return executed at the end of the interrupt service routine will then continue execution at the interrupted point. The 8 restart instructions and the address called by them are listed below:

Restart	Opcode	Adress Called
0	C7H	0000H
1	CFH	0008H
2	D7H	0010H
3	DFH	0018H
4	E7H	0020H
5	EFH	0028H
6	F7H	0030H
7	FFH	0038H

Restart 7 is interesting for a simple interrupt implementation, as will be seen later.

Since an interrupt can occur at any time (if enabled) the interrupt service routine must save any flags and registers it will use and restore them before returning. This is done with the PUSH and POP instructions. Interrupts should be reenabled before returning, because the processor disables further interrupts when it acknowledges one.

Since the Micropolis disk system is software controlled, it uses software timing loops in its operation. Any interrupts during one of these loops or during a disk read or write can cause errors. For this reason, the software disk driver disables interrupts when it is called, preventing any interruptions in its execution. Interupts must be reenabled at the completion of the sector data transfer when using interrupts in the system. Space is provided for placing an EI (FBH) instruction in the RES module. This byte, normally a NOP (00H) is located at 0A04H in version 4 of the Micropolis PDS. Any interrupt software initialization routines must change this location to the enable interrupt instruction so the disk driver will reenable interrupts at the completion of the sector transfer.

#### REAL TIME INTERRUPT HARDWARE

This circuit divides the 2 MHZ bus clock down to 1/60 hz (1 min period) and uses this divided rate to generate an interrupt. 7 divide-by-10 counters and one divide-by-12 counter are used to get the required divide by 1.2 \* 108.

Counters Al through A7 consist of two sections, a divide by two and a divide by five. The output of the divide by two (pin 12) is fed into the divide by five (pin 1). A8 is similar except the second section is a divide by six instead of a divide by

five. The reset to zero inputs on Al through A8 are grounded and the reset to nine inputs on Al through A7 are grounded to allow the counters to count normally.

The output of the counter chain (A8 pin 8) feeds a D-latch which latches the interrupt request.

A rising edge on the counter output (A8 pin 8) will clock the "1" on the D input (A9 pin 3) to the Q output (A9 pin 5). The Q output of the latch drives the input of an inverter (A10 pin 1) whose output drives the PINT (S100 pin 73) line. This line is active level low so the inverter provides the proper signal polarity as well as isolating the latch output from the bus.

An open collector inverter is used to drive this line as it is pulled up on the processor card in most systems. This allows any other possible interrupts to be OR tied together. The software must determine the cause of the interrupt if multiple interrupts are used this way.

When the CPU acknowledges the interrupt, SINTA is asserted on S-100 pin 96 which will clear the interrupt. This is inverted to provide the proper polarity for the clear input on the interrupt latch.

The next rising edge on the input of the D-latch will then generate another interrupt request 1 minute later.

Systems implementing a vectored interrupt scheme may use one of the lines VIO to VI7 (S-100 pins 4-11) instead of PINT. The vectored interrupt hardware will then generate the appropriate restart instruction on the interrupt acknowledge cycle. The service routine for the real time clock must then have a jump instruction at the appropriate restart vector location.

If multiple interrupts are employed in the system, a more sophisticated means of reseting the real time clock interrupt must be employed. This is to prevent the interrupt acknowledge cycle from another interrupt from clearing a pending clock interrupt. This can be accomplished by some additional hardware and using an OUT instruction to clear the interrupt in the clock service routine.

The input to the counter chain is taken from the clock line, S-100 pin 49. The required frequency is 2 Mhz for this counter. This line was used rather than phase 2 because phase 2 (pin 24) is at the processor speed, which is 4 MHZ in some systems. Clock, pin 49 is specified (by the IEEE) to always be 2 MHZ, independent of processor speed. This should be verified, as some older CPU cards put out a clock at the processor speed rather than 2 MHZ.

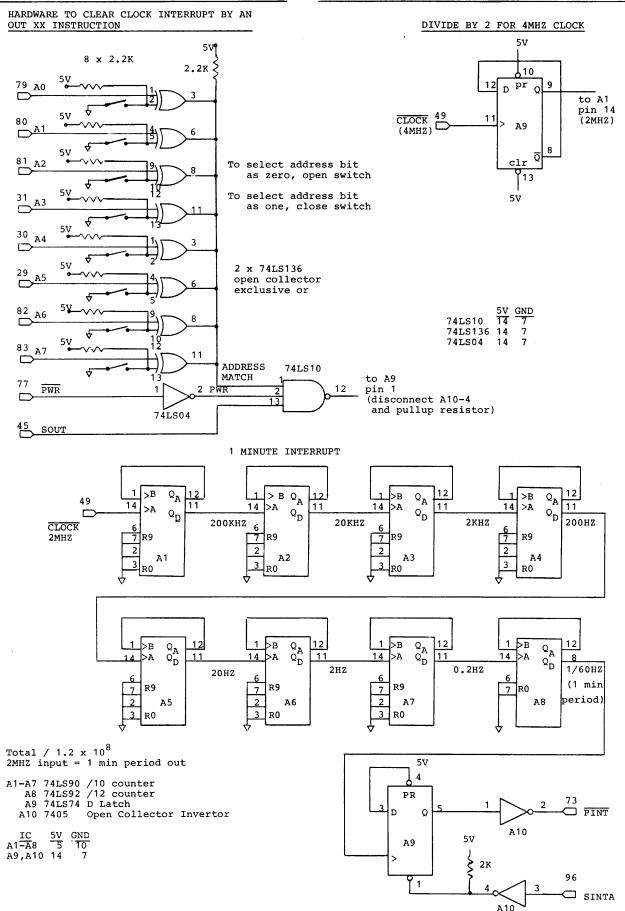
In my system, this is the only interrupt used so nothing drives the data in bus during the interrupt acknowledge cycle. The processor therfore reads an OFFH which is a Restart-7, which is the interrupt vector I used in the Real Time program. The data in bus must be pulled up to 5 volts through resisters to assure it will read as all ones when it is undriven.

This allows a simple implementation when a single interrupt is used in the system.

#### INDEX TO MUG'S FIRST YEAR

Ken Findlay (937 Briar Hill Ave., Toronto, Ontario M6B 1M1) has researched all of the MUG's first 12 issues, and coded the information by program name/topic author, program/article type, and category. He then wrote a set of programs to input the information to a data base, to manipulate it, and to print it. Although the printout can be in many forms, what follows is an index in alphabetical order by programname/topic.

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## MICROPOLIS USER'S GROUP NEWSLETTER INDEX ISSUES 1 THROUGH 12

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PGMNAME/TOPIC	CO./AUTHOR	PGM/ARTICLE TYPE	CATEGORY	VOL-PG	PGMNAME/TOPIC	CO./AUTHOR	PGM/ARTICLE TYPE	CATEGORY	VOL-PG	NEWS
		HIGH LEVEL LANGUAGE	COMPILER	012-03	HAM PGMS AVAILABLE		8080 APPL PGMS	HAM RADIO	005-12	LE
AMORT	MUG LIBR	BASIC APPL PGM	BUSINESS	008-05	HIGH LEVEL LANGUAGES		GENERAL INFO	IAH NADIO	005-11	ETTER
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BSS CCA DMS	CUSTOM ELEC	BASIC APPL PGM BASIC APPL PGM	BUSINESS DATA BASE	008-03	MICROPOLIS HARD DISKS	MICROPOLIS	HARDWARE	REFERENCE	005-05	
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DATABASE	J.SHAPIRO	BASIC APPL PGM	DATA BASE	003-07	MUG NEWSLETTER		GENERAL INFO		009-01	
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oldin III CO	0.001111	BASIC FOM IECHNIQUE		011-01	SAVING VARIADIES		BASIC UTILITY PGM		010-02	17

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PGMNAME/TOPIC	CO./AUTHOR	PGM/ARTICLE TYPE	CATEGORY	VOL-PG
SCREEN DISPLAY TEST*	J. CALLAWAY	BASIC UTILITY PGM	TERM I/O	006-06
SCREEN DISPLAY TEST*	B.RUDOW	BASIC UTILITY PGM	TERM I/O	006-07
SIZES		BASIC STATEMENT	REFERENCE	002-03
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YES/NO INPUT RESPONSES*	ED BURKHARDT	BASIC PGM TECHNIQUE		012-05

The programs and data file will be put on Library Disk 7 - which isn't done yet. You'll be able to use the programs for indexing magazine articles, records, books - whatever. They are a very useful set of routines. . . . . . . . .

#### LIBRARY DISKS 4 & 5??

I have succeeded in thoughly confusing a lot of people last month by referring to library disks 4 and 5. What I did, without telling you, was to make the Membership Directory be Disk 4, and the S/W Vendor's Directory be Disk 5.

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WANTED: People to input names and addresses at five cents each. I supply the programs, the disk, and the data. You mail the completed input back to me. Estimated need is for the first quarter of 1982. Tim Dawalt, P.O. Box 253, Lightstreet PA 17839, 717/784-4496

WANTED: A person to perform a disk conversion from 5 1/4 Micropolis CP/M to 8" Standard (and/or double-density) CP/M. Martin Rothstein, 21 E. 40 St., NY NY 10016, 212/683-5310

NOTICE: Micropolis will present a floppy disk maintenance course on Oct. 7 & 8 at the Los Angeles Airport Marriott. Cost is \$200 for first student, \$175 for additional students from the same company. Fees include lunches, manuals and diagnostic software. Contact Bob Louch at 213/709-3300.

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